

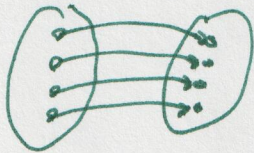
A.

(1)

$f: A \rightarrow B$  is a bijection if it is one-to-one and onto  
(injective) (surjective)

- one-to-one: for each  $y \in B$  there is at most one  $x \in A$  s.t.  $f(x) = y$ .

- onto: for each  $y \in B$  there is at least one  $x \in A$  s.t.  $f(x) = y$ .



Bijection Rule: If there exists a bijection  $f: A \rightarrow B$  then  $|A| = |B|$

Product Rule:

A = "the set of possible executions of a procedure P"

B = "the set we want to know the size of"

X = "an n-element set"

$$X = \{a, b, c\}$$

$P(X)$  = the set of all subsets of X

$$P(X) = \left\{ \begin{array}{l} \emptyset, \{a\}, \{b\}, \{c\} \\ \{a, b\}, \{a, c\}, \{b, c\}, \\ \{a, b, c\} \end{array} \right\}$$

Theorem: For any n-element set X,  $|P(X)| = 2^n$ .

Proof: Let  $X = \{x_1, \dots, x_n\}$ .

For any  $S \subseteq X$ , define the bitstring  ~~$b_1 \dots b_n$~~  where

$$f(S) = b_1 \dots b_n \text{ where } b_i = \begin{cases} 0 & \text{if } x_i \notin S \\ 1 & \text{if } x_i \in S. \end{cases}$$

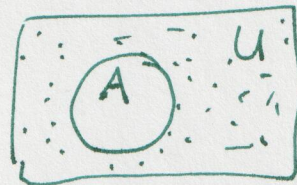
Then  $f$  is a bijection from  $P(X)$  onto binary strings of length  $n$ .

Can also prove using product rule.

Complement Rule: If  $A \subseteq U$  then  $|A| = |U| - |U \setminus A|$

(2)

Example: Password is a string of 8 characters over the alphabet  $\Sigma = \{a, b, \dots, 0, 1, 2, \dots, 9\}$  that contains at least one digit.



$S =$  "the set of all passwords"

$U =$  "all 8-character strings over  $\Sigma$ "

$|\Sigma| = 36$

$|U| = 36^8$  (Product Rule).

~~Let~~

$U \setminus S =$  "all 8 character strings with no digits."  
= "8 character strings over  $\{a, b, c, \dots, z\}$ ."

$|U \setminus S| = 26^8$

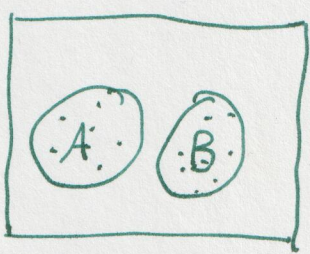
$\therefore |S| = |U| - |U \setminus S| = 36^8 - 26^8 = 2,612,282,842,880$

Not resistant to brute force.

# The Sum Rule.

If A and B are disjoint, then  $|A \cup B| = |A| + |B|$

If  $A_1, \dots, A_n$  are pairwise disjoint then  $|A_1 \cup \dots \cup A_n| = |A_1| + |A_2| + \dots + |A_n|$ .



Example: A valid password is any string of length 6, 7, or 8

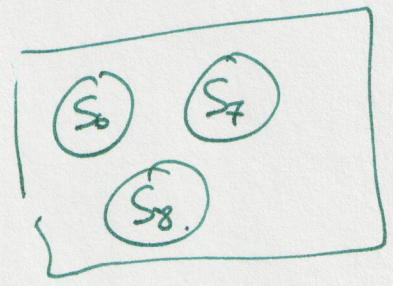
over the alphabet  $\Sigma = \{a, \dots, z, 0, \dots, 9\}$  where that contains at least one character.

For each  $n \in \{6, 7, 8\}$ , let  $S_n$  be the set of valid passwords of length  $n$ .

$|S_6| + |S_7| + |S_8|$

$$S = S_6 \cup S_7 \cup S_8 \Rightarrow |S| = |S_6 \cup S_7 \cup S_8| = 36^6 - 26^6 + 36^7 - 26^7 + 36^8 - 26^8$$

$$= 2,684,483,063,360.$$



P. I-E (3 sets).

For any sets A, B, and C  $|A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C|$ . ←

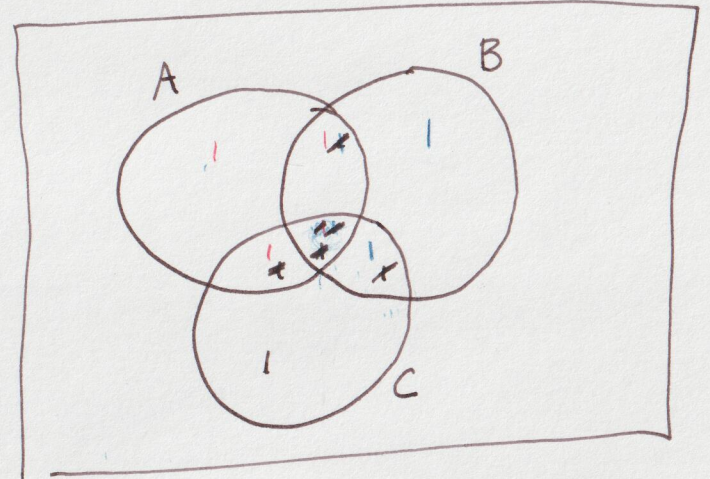
- same problem (bit strings).

C = length-17 bitstrings with 01 at positions 6 & 7.

$$|C| = 2^{15}$$

$$C = \underbrace{\text{LLLLL}}_5 01 \underbrace{\text{LLLLLLLLLLLLL}}_{10}$$

$$|A \cup B \cup C| = 2^{14} + 2^{15} + 2^{15} - 2^{12} - 2^{12} - 2^{13} + 2^{10} = 66,560$$



$$|A \cap C| = \underbrace{111}_{2} \underbrace{\text{LLL}}_3 01 \underbrace{\text{LLLLLLLLL}}_{10}$$

$$|B \cap C| = \underbrace{\text{LLLLL}}_5 01 \underbrace{\text{LLLLLLLLL}}_8 00$$

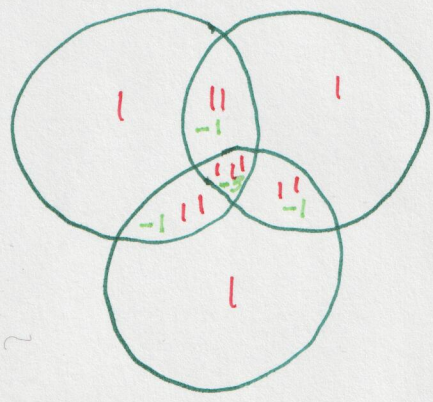
$$|A \cap B \cap C| = \underbrace{111}_{2} \underbrace{\text{LLL}}_3 01 \underbrace{\text{LLLLLLLLL}}_8 00$$

Principle Inc-Exc For 3 sets.

For any 3-sets A, B, C

$$|A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C|$$

General formula for n-terms has  $2^n - 1$  terms.



Example: ~~bit~~ Length-17 bitstrings again or have 10 at positions 7, 8.

$$|A| = 2^{14}$$

$$|B| = 2^{15}$$

$$|C| = 2^{15}$$

$$|A \cap B| = 2^{12}$$

$$|B \cap C| = 2^{13}$$

$$|A \cap C| = 2^{12}$$

$$|A \cap B \cap C| = 2^{10}$$

$$|A \cup B \cup C| = 2^{14} + 2^{15} + 2^{15} - 2^{12} - 2^{13} - 2^{12} + 2^{10} = 66,560.$$